

Data-Parallel Computing Meets STRIPS

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Outline

1 Motivation

2 DPPS

3 DPPS as Planning

Data Processing — Before “Big Data”

- Database Management Systems (DBMS)
- Declarative query — expressed in SQL
- Query execution plan
 - Easy to generate from declarative query
 - Hard to optimize
- Very limited support for user-defined functions

Data Processing — After “Big Data”

- MapReduce / Hadoop / Dryad
 - Low-level programming
 - Only user-defined functions
 - No declarative queries
- SCOPE / DryadLINQ / Pig / Hive
 - High-level programming
 - Support user-defined functions
 - Limited declarative queries

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User Defined Functions in Declarative Queries

- Including user-defined functions hinders query optimization
 - User must specify some base plan
 - Query plan optimizer does not “understand” user-defined functions, and does not know which optimizations are safe
- Existing approaches:
 - No optimization when user-defined function in query
 - User-defined functions must have some pre-specified signature
 - Static code analysis to “understand” user-defined functions

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Running Example: Histogram Computation

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Query Execution Plan

```
Agg (age, scan(T))  
Agg (rls, scan(T))
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Running Example: Histogram Computation (2)

- Suppose we have a user-defined function, DAgg, which aggregates by two fields simultaneously
- The question is how to come up with this execution plan automatically

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Query Execution Plan using DAgg

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Our Contribution

- Introduce Data-Parallel Program Synthesis (DPPS), a formal framework for studying these problems
- Study expressivity and complexity of DPPS
- Show compilation to AI planning

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Data-Parallel Program Synthesis Framework

- Framework is based on tracking *data chunks*
- A data chunk represents some piece of data, e.g.:
 - all records of males between the ages of 18–49
 - the average salary of all males between the ages of 18–49
- We do not need to know the *value* of the data, only its description
- Each data chunk d is associated with the amount σ_d of memory it requires

DPPS Task

- D — a set of possible data chunks, with sizes σ_d
- N — a finite set of computing units, with memory capacities κ_n
- A — a set of possible computation primitives, $a \in A$ described by:
 - $\bar{I} \subseteq D$ is the required input
 - $\bar{O} \subseteq D$ is the produced output
 - $C : N \rightarrow \mathbb{R}^{0+}$ computation cost on each processor
- $T : N \times D \times N \rightarrow \mathbb{R}^{0+}$ — the data transmission cost function
- s_0 — the initial state of the computation
- G — the goal of the computation

DPPS Task (2)

- A DPPS state specifies which processor holds which data chunks
- A solution is a sequence of actions (compute / transmit / delete data) which achieves the goal from the initial state
- The possible data chunks D and computations A may be given explicitly or described implicitly
 - If they are described implicitly the sets could be infinite

DPPS Expressivity

Theorem

DPPS is at least as expressive as relational algebra with aggregation

Proof sketch.

Given a relational algebra expression, we can construct a DPPS task whose operators are the RA operators, and data chunks are possible RA expressions. □

DPPS Complexity ☺

Theorem

Satisficing data-parallel program synthesis is NP-hard, even when the possible data chunks are given explicitly.

Proof sketch.

By reduction from SAT, exploiting memory capacity constraints



DPPS Complexity ☹☹

Theorem

Optimal data-parallel program synthesis with a single processor is NP-hard, even if the possible data chunks are given explicitly, and there are no memory constraints.

Proof sketch.

By reduction from delete-free planning



DPPS Complexity ☹☹☹

Theorem

Optimal data-parallel program synthesis with a single data chunk is NP-hard.

Proof sketch.

By reduction from the Steiner tree problem



DPPS Complexity ☺

Theorem

Satisficing data-parallel program synthesis with no memory constraints can be solved in polynomial time, when the possible data chunks are given explicitly.

Proof sketch.

By reduction from delete-free planning



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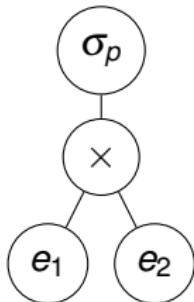
DPPS Compilation

- When the computations and data chunks are given explicitly, compilation to planning is straightforward
 - Predicate `holds` (`?node, ?data`)
 - Actions
 - For each computation `compute` (`?node, ?computation`)
 - Transmission `transmit` (`?node, ?data, ?node2`)
 - Data deletion `del` (`?node, ?data`)
 - Capacity constraints can be enforced with numerical fluents

DPPS Compilation without Explicit Data

- When the computations and data chunks are given implicitly, compilation is still possible sometimes
- When data chunks have a structure (e.g., expression trees), it is possible to represent such trees using predicates

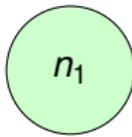
Expression Tree



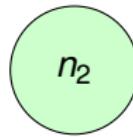
Encoding

```
select (n1, p, n2)
join (n2, e1, e2)
```

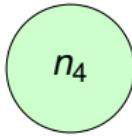
DPPS Compilation: Proof of Concept



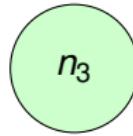
$\sigma_{hash(PK)=1}(T)$



$\sigma_{hash(PK)=2}(T)$

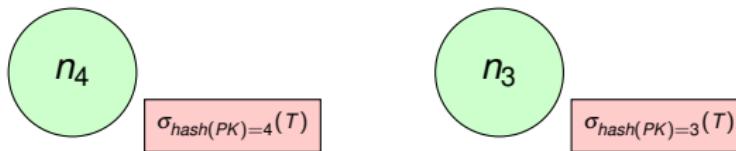
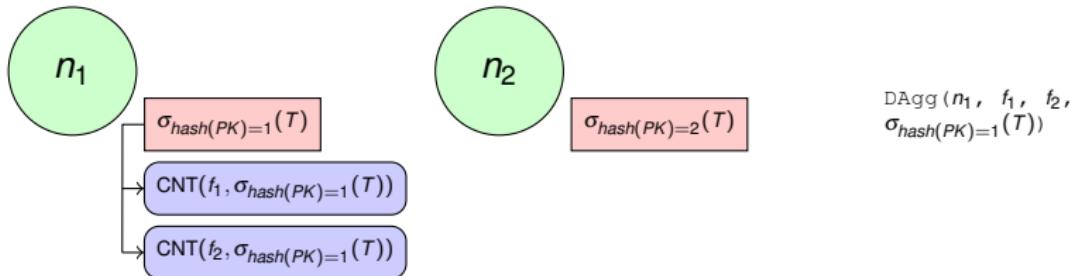


$\sigma_{hash(PK)=4}(T)$

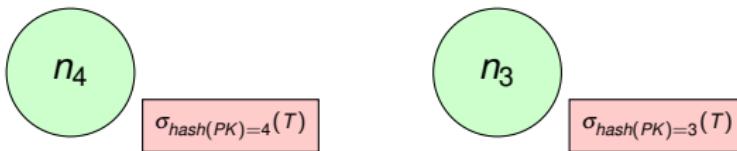
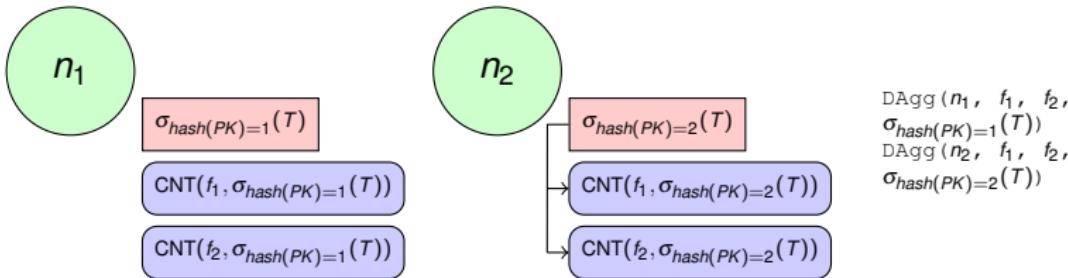


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DPPS Compilation: Proof of Concept

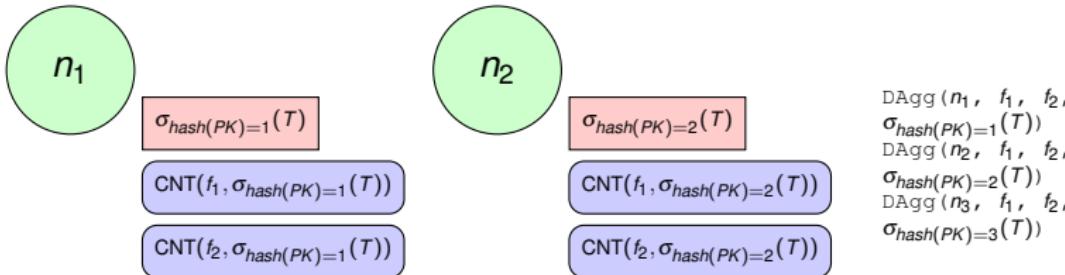


DPPS Compilation: Proof of Concept

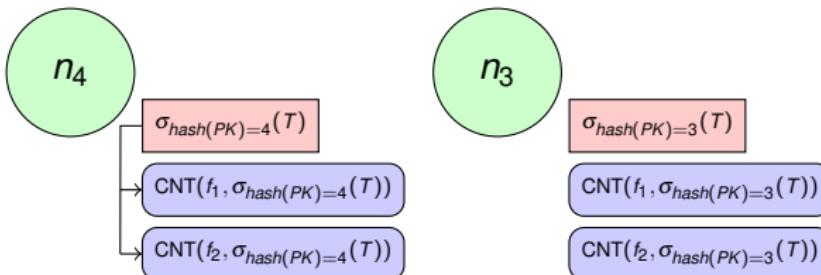
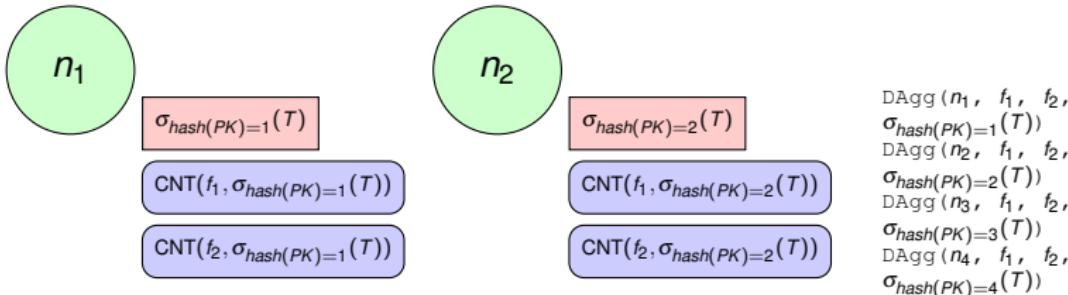


DPPS Compilation: Proof of Concept

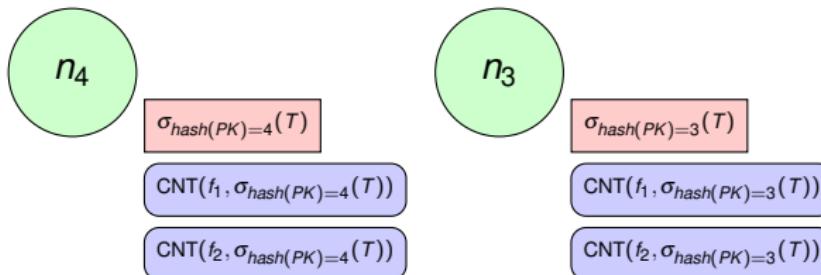
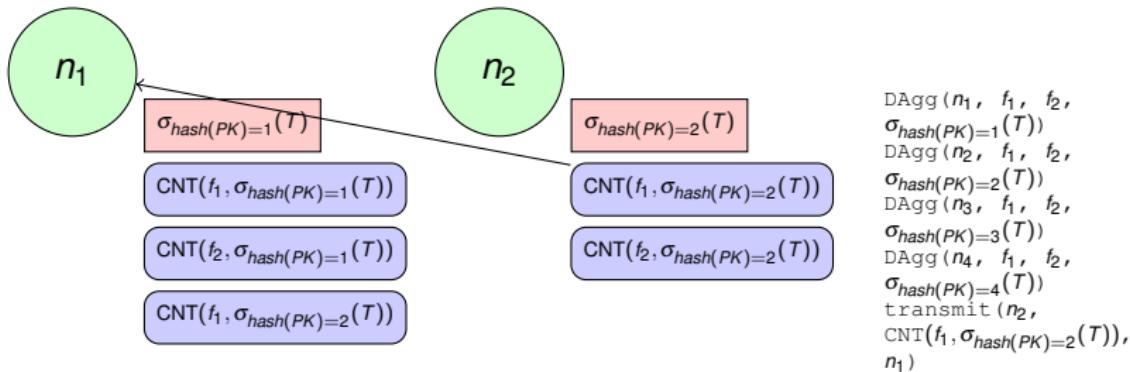
Diagram illustrating DPPS Compilation for four nodes (n_1, n_2, n_3, n_4) showing their local states and aggregated results.



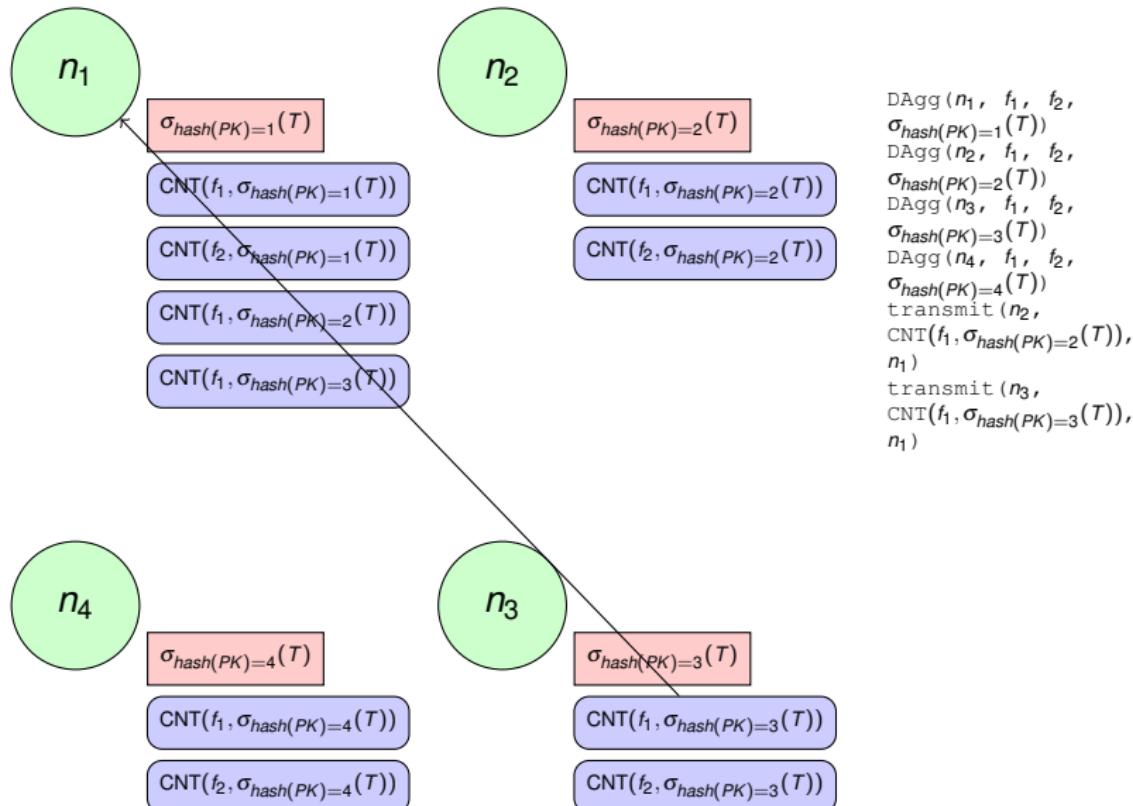
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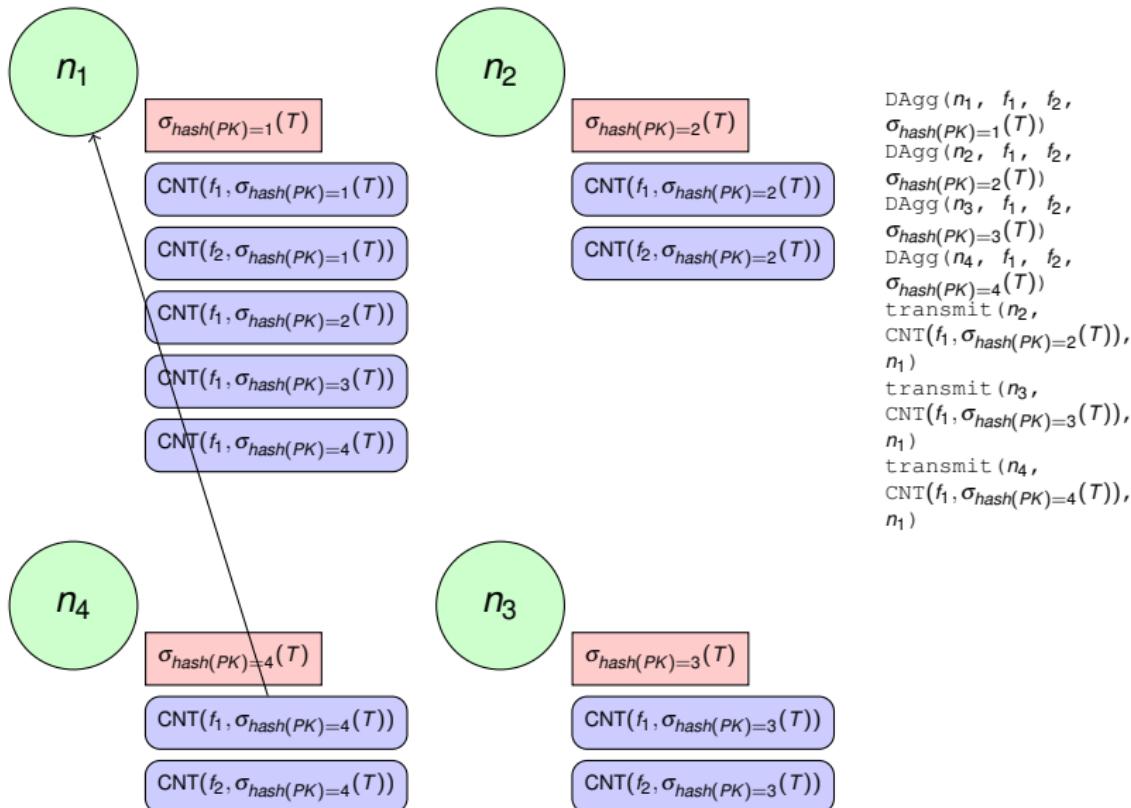
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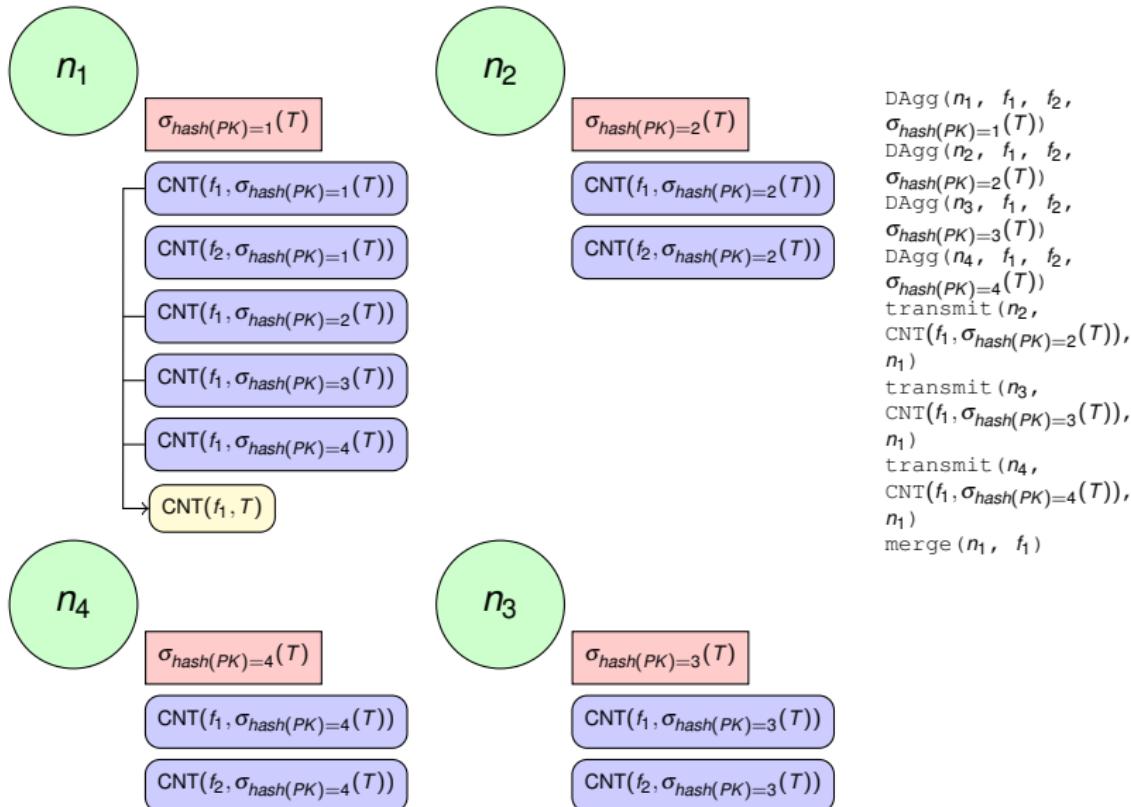
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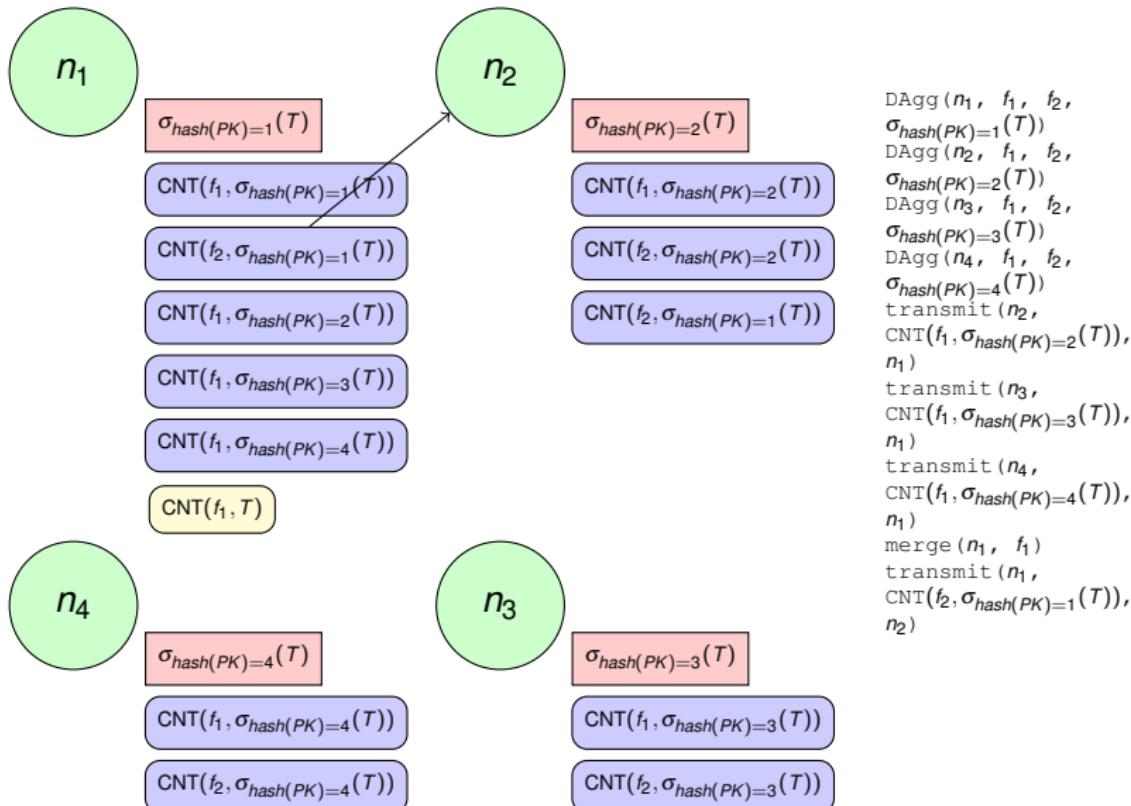
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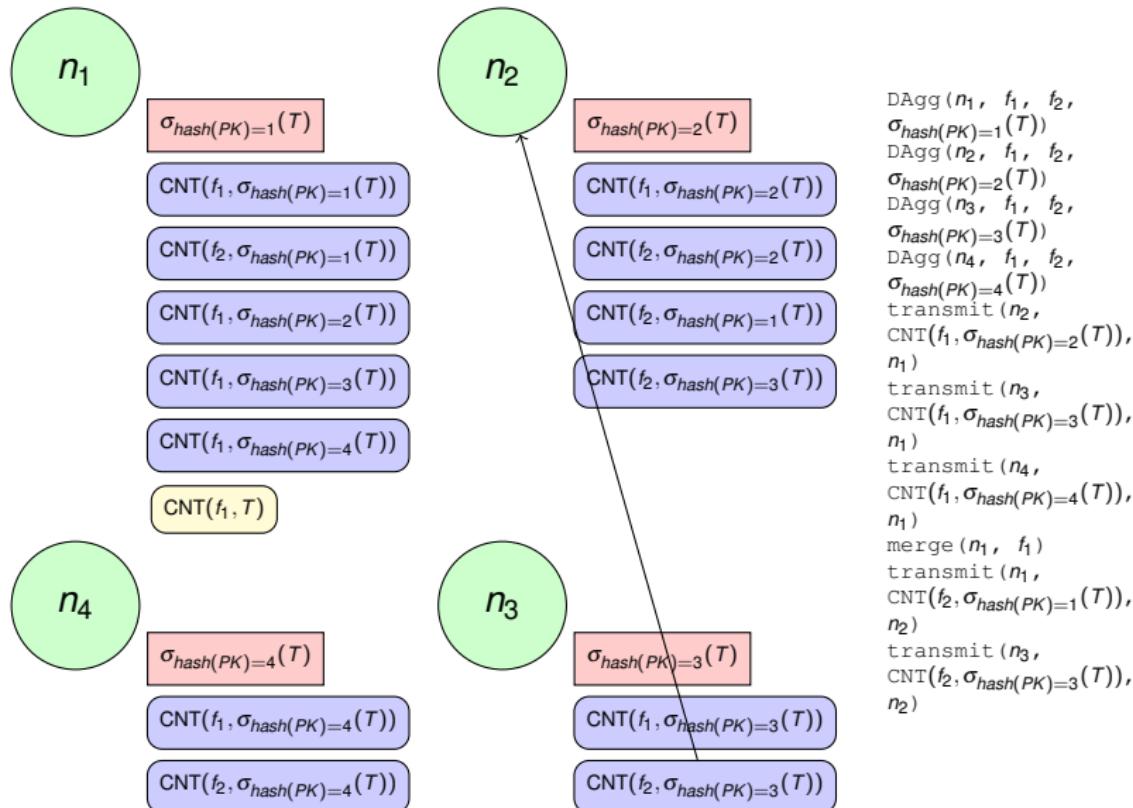
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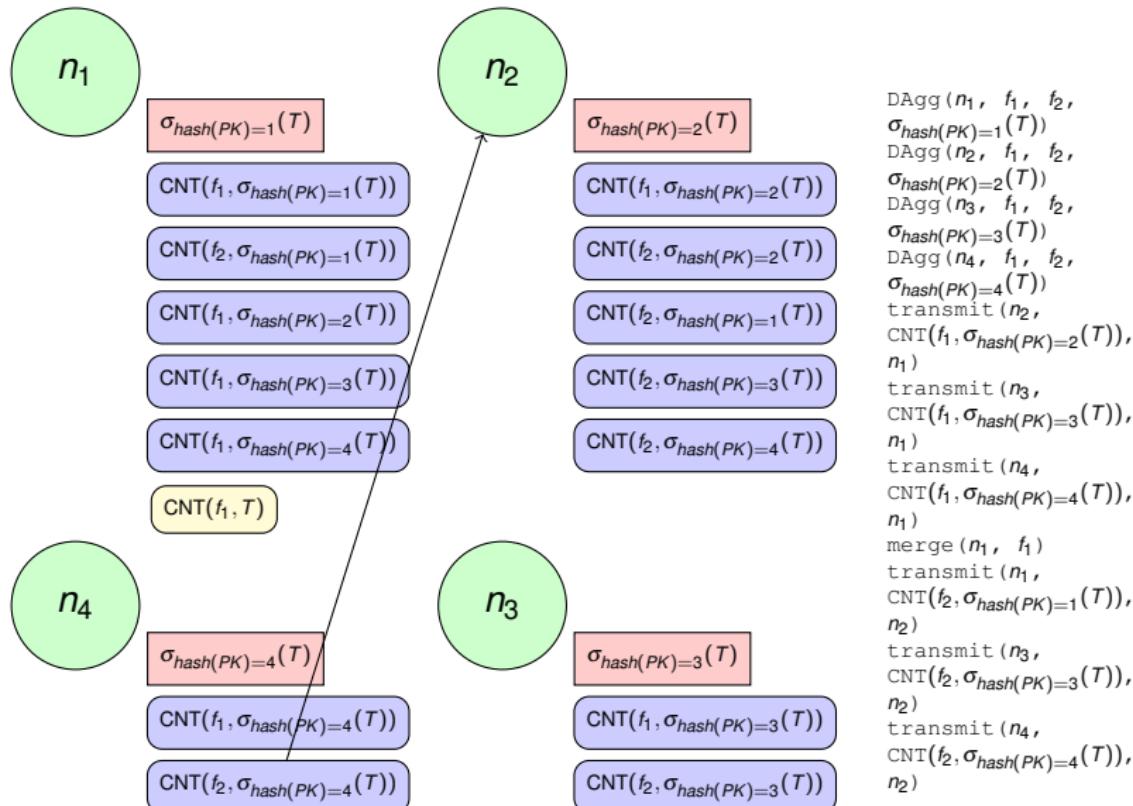
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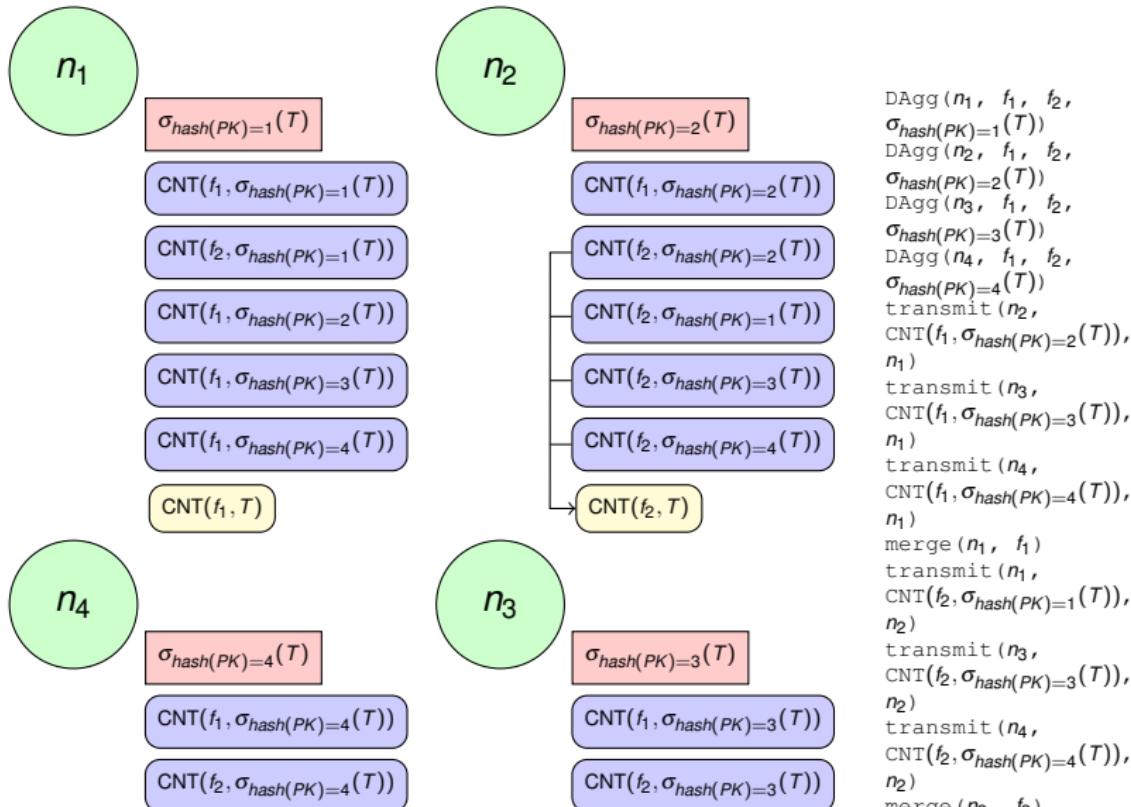
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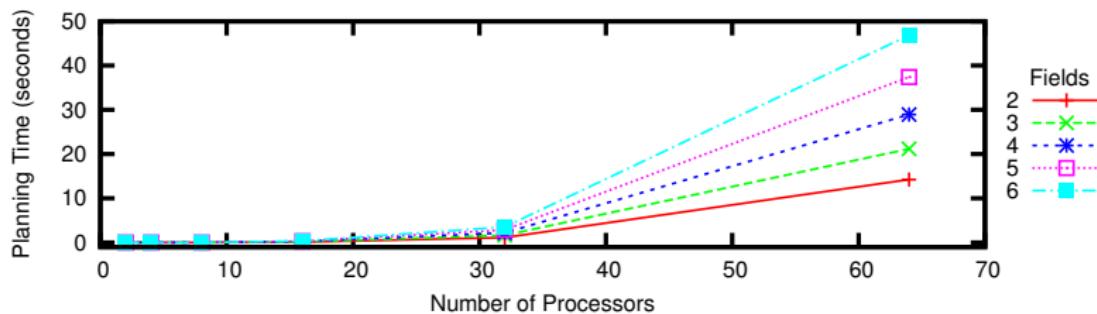
DPPS Compilation: Proof of Concept



DPPS Compilation: Proof of Concept



DPPS Compilation: Proof of Concept



- Histogram of F fields of a table divided across N processors
- Solved by GBFS using relaxed plan heuristic in Fast Downward
- Solutions were optimal (although this is not guaranteed)

Summary

- DPPS is a flexible framework for describing data-parallel computations
- Solving DPPS is possible through compilation to AI planning
- We expect DPPS to lead to interesting questions in AI planning

Thank You